

Self-Indifferentiability of $C(X) := H(X)$ from a RO: $n \rightarrow n$ where H is RO: $n \rightarrow n$.

G0	G1	G2	G3	G4	G5	G6	G7
C(X): $Y = H(X)$ Return Y	C(X): If $(X,Y) \notin LC$ $Y = H(X)$ $LC \leftarrow (X,Y) : LC$ Else Find Y Return Y	C(X): If $(X,Y) \notin LC$ If $(X,Y) \notin L$ $Y \leftarrow \{0,1\}^n$ $L \leftarrow (X,Y) : L$ Else Find Y $LC \leftarrow (X,Y) : LC$ Else Find Y Return Y	Unmodified	C(X): If $(X,Y) \notin LC$ If $(X,Y) \notin L$ $Y \leftarrow \{0,1\}^n$ $// L \leftarrow (X,Y) : L$ Else Find Y $LC \leftarrow (X,Y) : LC$ Else Find Y Return Y	Unmodified	C(X): If $(X,Y) \notin LC$ $// If (X,Y) \notin L$ $Y \leftarrow \{0,1\}^n$ $// Else Find Y$ $LC \leftarrow (X,Y) : LC$ Else Find Y Return Y	C(X): If $(X,Y) \notin LC$ $Y \leftarrow \{0,1\}^n$ $LC \leftarrow (X,Y) : LC$ Else Find Y Return Y
H(X): If $(X,Y) \notin L$ $Y \leftarrow \{0,1\}^n$ $L \leftarrow (X,Y) : L$ Else Find Y Return Y	Unmodified	Unmodified	H(X): If $(X,Y) \notin L$ If $(X,Y) \in LC$ $L \leftarrow (X,Y) : L$ Else $Y \leftarrow \{0,1\}^n$ $L \leftarrow (X,Y) : L$ Else Find Y Return Y	Unmodified	H(X): If $(X,Y) \notin L$ If $(X,Y) \in LC$ $L \leftarrow (X,Y) : L$ Else $Y \leftarrow \{0,1\}^n$ $L \leftarrow (X,Y) : L$ $LC \leftarrow (X,Y) : LC$ Else Find Y Return Y	Unmodified	H(X): If $(X,Y) \notin L$ $Y \leftarrow C(X)$ $L \leftarrow (X,Y) : L$ Else Find Y Return Y

G0: The real game.

G1: Keep track of C queries.

G2: Inline the query to H within C.

G3: Prepare for moving handling of list L from C to H.

Add all the relevant entries in LC to L. (At this point this change won't add anything yet.)

G4: Remove handling of L from C.

An inconsistency could arise if X was queried to C (and was not added to L under C) and then later queried to H. But in this case H2 updates L with entries in LC.

G5: Prepare for moving checking list L from C.

Add all entries in L to LC. Preemptively complete chains.

G6: Anything in L is on LC so the check in the If statement always passed (and Else is never executed).

G7: Instead of LC entries use oracle C.

If the entry was on LC, C would return the same entry as in G6.

If not, in both G6 and G7 a independent entry is sampled and added to LC. In both games this new sample is returned under H.

The resulting simulator is:

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Sim^C(X):
Return C(X)
    
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Range Extension: Indifferentiability of $C(X) := H(0|X) | H(1|X)$ from a RO: $n-1 \rightarrow 2n$ where H is RO: $n \rightarrow n$.

G0	G1	G2	G3	G4	G5	G6	G7	G8	G9
C(X): $Y0 = H(0 X)$ $Y1 = H(1 X)$ Return $Y0 Y1$	C(X): If $(X,Y0 Y1) \notin LC$ $Y0 = H(0 X)$ $Y1 = H(1 X)$ LC $\leftarrow (X,Y1 Y2):LC$ Else Find $Y0 Y1$ in LC Return $Y0 Y1$	Unmodified	C(X): If $(X,Y0 Y1) \notin LC$ If $(0 X,Y) \notin L$ $Y0,Y1 \leftarrow \{0,1\}^n$ $L \leftarrow (0 X,Y0) : L$ $L \leftarrow (1 X,Y1) : L$ Else Find $Y0$ If $(1 X,Y) \notin L$ $Y1 \leftarrow \{0,1\}^n$ $L \leftarrow (1 X,Y1) : L$ Else Find $Y1$ LC $\leftarrow (X,Y0 Y1):LC$ Else Find $Y0 Y1$ in LC Return $Y0 Y1$	C(X): If $(X,Y0 Y1) \notin LC$ If $(0 X,Y) \notin L$ $Y0,Y1 \leftarrow \{0,1\}^n$ $L \leftarrow (0 X,Y0) : L$ $L \leftarrow (1 X,Y1) : L$ Else Find $Y0,Y1$ LC $\leftarrow (X,Y0 Y1):LC$ Else Find $Y0 Y1$ in LC Return $Y0 Y1$	Unmodified	C(X): If $(X,Y0 Y1) \notin LC$ If $(0 X,Y) \notin L$ $Y0,Y1 \leftarrow \{0,1\}^n$ // $L \leftarrow (0 X,Y0) : L$ // $L \leftarrow (1 X,Y1) : L$ Else Find $Y0,Y1$ LC $\leftarrow (X,Y0 Y1):LC$ Else Find $Y0 Y1$ in LC Return $Y0 Y1$	Unmodified	C(X): If $(X,Y0 Y1) \notin LC$ // If $(0 X,Y) \notin L$ $Y0,Y1 \leftarrow \{0,1\}^n$ // Else Find $Y0,Y1$ LC $\leftarrow (X,Y0 Y1) : LC$ Else Find $Y0 Y1$ in LC Return $Y0 Y1$	Unmodified
H(b X): If $(b X,Yb) \notin L$ $Yb \leftarrow \{0,1\}^n$ $L \leftarrow (b X,Yb) : L$ Else Find Yb in L Return Yb	Unmodified	H(b X): If $(b X,Yb) \notin L$ $Y0,Y1 \leftarrow \{0,1\}^n$ $L \leftarrow (0 X,Y0) : L$ $L \leftarrow (1 X,Y1) : L$ Else Find Yb in L Return Yb	Unmodified	Unmodified	H(b X): If $(b X,Yb) \notin L$ If $(X,Y0 Y1) \in LC$ $L \leftarrow (0 X,Y0) : L$ $L \leftarrow (1 X,Y1) : L$ Else $Y0,Y1 \leftarrow \{0,1\}^n$ $L \leftarrow (0 X,Y0) : L$ $L \leftarrow (1 X,Y1) : L$ Else Find Yb in L Return Yb	Unmodified	H(b X): If $(b X,Yb) \notin L$ If $(X,Y0 Y1) \in LC$ $L \leftarrow (0 X,Y0) : L$ $L \leftarrow (1 X,Y1) : L$ Else $Y0,Y1 \leftarrow \{0,1\}^n$ $L \leftarrow (0 X,Y0) : L$ $L \leftarrow (1 X,Y1) : L$ LC $\leftarrow (X,Y0 Y1):LC$ Else Find Yb in L Return Yb	Unmodified	H(b X): If $(b X,Yb) \notin L$ $Y0 Y1 \leftarrow C(X)$ $L \leftarrow (0 X,Y0) : L$ $L \leftarrow (1 X,Y1) : L$ Else Find Yb in L Return Yb

G0: The real game.

G1: Keep track of C queries.

G2: Add entries for $b = 0$ and $b = 1$ in tandem: For each b we have that $(b|X,Yb) \notin L$ iff $(0|X,Y0) \notin L$ and $(1|X,Y1) \notin L$.

G3: Inline the two H queries within C.

G4: Simplify due to change in G2.

G5: Prepare for moving handling of list L from C to H. For all entries X in LC add $0|X$ and $1|X$ to L. (This does nothing at this step since for X on LC has $0|X$ and $1|X$ are both in L.)

G6: Remove handling of L from C. H will handle the entries consistently since it looks at entries on LC.

G7: Always add the L entry to LC. Prepares us to stop relying on LC and use C instead.

G8: If some entry is in L it will be on LC too (due to G7). So remove If-Else.

G9: Use C instead of LC entries irrespectively of whether or not (X,Y) is on LC.

Simplifying H, we get the simulator below, which makes 1 C-query per invocation:

Sim^C(b|X):
 $Y0|Y1 \leftarrow C(X)$
 Return Yb

Question: Is a this range extender *multi-stage* indifferentiable?

Answer: Yes, because the simulator is stateless!

Domain Extension: Indifferentiability of $C(X) := H_2(H_1(X_1, X_2), X_3)$ from a RO: $3n \rightarrow n$ where H_1, H_2 are independent RO: $2n \rightarrow n$.

G0	G1	G2	G3	G4	G5
<p><u>C(X1,X2,X3):</u> $Y1 = H1(X1, X2)$ $Y2 = H2(Y1, X3)$ Return Y2</p>	<p><u>C(X1,X2,X3):</u> If $((X1, X2, X3), Y2) \notin LC$ $Y1 = H1(X1, X2)$ $Y2 = H2(Y1, X3)$ $LC \leftarrow ((X1, X2, X3), Y2) : LC$ Else Find $((X1, X2, X3), Y2) \in LC$ Return Y2</p>	<p><u>C(X1,X2,X3):</u> If $((X1, X2, X3), Y2) \notin LC$ $Y1 = H1(X1, X2)$ If $((Y1, X3), Y2) \notin L2$ $Y2 \leftarrow \{0,1\}^n$ $L2 \leftarrow ((Y1, X3), Y2) : L2$ Else Find Y2 in L2 $LC \leftarrow ((X1, X2, X3), Y2) : LC$ Else Find $((X1, X2, X3), Y2) \in LC$ Return Y2</p>	<p>Unmodified</p>	<p><u>C(X1,X2,X3):</u> If $((X1, X2, X3), Y2) \notin LC$ $Y1 = H1(X1, X2)$ If $((Y1, X3), Y2) \notin L2$ $Y2 \leftarrow \{0,1\}^n$ $// L2 \leftarrow ((Y1, X3), Y2) : L2$ Else Find Y2 in L2 $LC \leftarrow ((X1, X2, X3), Y2) : LC$ Else Find $((X1, X2, X3), Y2) \in LC$ Return Y2</p>	<p>Unmodified</p>
<p><u>H1(X1,X2):</u> If $((X1, X2), Y1) \notin L1$ $Y1 \leftarrow \{0,1\}^n$ $L1 \leftarrow ((X1, X2), Y1) : L1$ Else Find Y1 in L1 Return Y1</p>	<p>Unmodified</p>	<p>Unmodified</p>	<p><u>H2(Y1,X3):</u> If $((Y1, X3), Y2) \notin L2$: If $((X1, X2, X3), Y2) \in LC \ \& \ ((X1, X2), Y1) \in L1$ $L2 \leftarrow ((Y1, X3), Y2) : L2$ Else $Y2 \leftarrow \{0,1\}^n$ $L2 \leftarrow ((Y1, X3), Y2) : L2$ Else Find Y2 in L2 Return Y2</p>	<p>Unmodified</p>	<p><u>H2(Y1,X3):</u> If $((Y1, X3), Y2) \notin L2$: If $((X1, X2, X3), Y2) \in LC \ \& \ ((X1, X2), Y1) \in L1$ $L2 \leftarrow ((Y1, X3), Y2) : L2$ Else $Y2 \leftarrow \{0,1\}^n$ If $((X1, X2), Y1) \in L1$ $LC \leftarrow ((X1, X2, X3), Y2) : LC$ $L2 \leftarrow ((Y1, X3), Y2) : L2$ Else Find Y2 in L2 Return Y2</p>

G0: The real game.

G1: Keep track of C queries.

G2: Inline H2.

G3: Add C queries relevant to the current H2 query to L2. (Does not execute at this point.)

G4: Remove handling of L2 from C. OK due to change in G3 and keeping track of C queries introduced in G1.

G5: **Preemptively** create LC entry for H1 and H2 entries that form a chain.

G6	G7	G8	G9	G10	G11
Unmodified	C(X1,X2,X3): If ((X1,X2,X3),Y2) ∈ LC Y1 = H1(X1,X2) // If ((Y1,X3),Y2) ∈ L2 Y2 ← {0,1} ⁿ // Else Find Y2 in L2 LC ← ((X1,X2,X3),Y2) : LC Else Find ((X1,X2,X3),Y2) ∈ LC Return Y2	C(X1,X2,X3): If ((X1,X2,X3),Y2) ∈ LC // Y1 = H1(X1,X2) Y2 ← {0,1} ⁿ LC ← ((X1,X2,X3),Y2) : LC Else Find ((X1,X2,X3),Y2) ∈ LC Return Y2	Unmodified	Unmodified	Unmodified
H1(X1,X2): If ((X1,X2),Y1) ∈ L1 Y1 ← {0,1} ⁿ If ((Y1,*),*) ∈ L2 SetFlag L1 ← ((X1,X2),Y1) : L1 Else Find Y1 in L1 Return Y1	Unmodified	Unmodified	H1(X1,X2): If ((X1,X2),Y1) ∈ L1 Y1 ← {0,1} ⁿ If ((*,*),Y1) ∈ L1 SetFlag If ((Y1,*),*) ∈ L2 SetFlag L1 ← ((X1,X2),Y1) : L1 Else Find Y1 in L1 Return Y1	Unmodified	Unmodified
Unmodified	Unmodified	Unmodified	Unmodified	H2(Y1,X3): If ((Y1,X3),Y2) ∈ L2: If ((X1,X2,X3),Y2) ∈ LC & ((X1,X2),Y1) ∈ L1 L2 ← ((Y1,X3),Y2) : L2 Else If ((X1,X2),Y1) ∈ L1 Y2 ← {0,1} ⁿ LC ← ((X1,X2,X3),Y2) : LC Else Y2 ← {0,1} ⁿ L2 ← ((Y1,X3),Y2) : L2 Else Find Y2 in L2 Return Y2	H2(Y1,X3): If ((Y1,X3),Y2) ∈ L2: If ((X1,X2),Y1) ∈ L1 Y2 ← C(X1,X2,X3) L2 ← ((Y1,X3),Y2) : L2 Else Y2 ← {0,1} ⁿ L2 ← ((Y1,X3),Y2) : L2 Else Find Y2 in L2 Return Y2

G6: Avoid **accidental chains**: Chains with where H1 query comes **after** a matching H2 query are excluded. This introduces a term: $(q_C + q_1)(q_C + q_2) / 2^n$

G7: Assume $((Y1,X3),Y2)$ is not in L2 (in which case Y1 is not needed in C). Suppose not, i.e., that $((Y1,X3),Y2)$ is in L2. We have two cases:

$((X1,X2),Y1)$ was on L1 when $((Y1,X3),Y2)$ was added to L2. Then $(X1,X2,X3,Y2)$ would be on LC due to the change in **G5**.

$((X1,X2),Y1)$ was not on L1 when $((Y1,X3),Y2)$ was added to L2. Then it will be added later on since the H1 query under C will certainly add it. This however would set Flag.

G8: Remove the H1 call (Y1 is not used in C). The H1 entry will be created when needed. Difference: $(q_C + q_1)(q_C + q_2) / 2^n$ (flag is set more often in **G7** than in **G8**).

G9: Avoid collisions in H1. This means two different chains with different Y1 values would have the same C values iff they have the same H2 values. Difference: $(q_C + q_1)(q_C + q_1) / 2^n$

G10: Use **different** Y2 for all entries $((X1,X2),Y1)$ was on L1. No change due to avoiding collisions in **G8** (entry on L1 is unique).

G11: Use C instead of LC. When $((X1,X2),Y1)$ is on L1 we either retrieve Y2 from LC or create a new one and added to LC. This is equivalent to a C call.

The resulting simulator, which makes at most 1 C-query per H2 call, is:

Sim1(X1,X2): If ((X1,X2),Y1) ∈ L1 Y1 ← {0,1} ⁿ If ((*,*),Y1) ∈ L1 SetFlag If ((Y1,*),*) ∈ L2 SetFlag L1 ← ((X1,X2),Y1) : L1 Else Find Y1 in L1 Return Y1	Sim2^C(Y1,X3): If ((Y1,X3),Y2) ∈ L2: If ((X1,X2),Y1) ∈ L1 Y2 ← C(X1,X2,X3) L2 ← ((Y1,X3),Y2) : L2 Else Y2 ← {0,1} ⁿ L2 ← ((Y1,X3),Y2) : L2 Else Find Y2 in L2 Return Y2
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Extra Example: Indifferentiability of $C(X) := H1(X) + H2(X)$ from a RO : $n \rightarrow n$ where $H1$ and $H2$ are independent RO: $n \rightarrow n$.

G0 C(X): $Y1 = H1(X)$ $Y2 = H2(X)$ $Z = Y1 + Y2$ Return Z	G1 C(X): If (X,Z) \notin LC $Y1 = H1(X)$ $Y2 = H2(X)$ $Z = Y1 + Y2$ LC \leftarrow (X,Z) : LC Else Find Z in LC Return Z	G2 C(X): If (X,Z) \notin LC If (X,Y1) \notin L1 $Y1 \leftarrow \{0,1\}^n$ L1 \leftarrow (X,Y1) : L1 Else Find Y1 $Y2 = H2(X)$ $Z = Y1 + Y2$ LC \leftarrow (X,Z) : LC Else Find Z in LC Return Z	G3 C(X): If (X,Z) \notin LC Y2 = H2(X) If (X,Y1) \notin L1 $Z \leftarrow \{0,1\}^n$ Y1 = Z + Y2 L1 \leftarrow (X,Y1) : L1 Else Find Y1 $Z = Y1 + Y2$ LC \leftarrow (X,Z) : LC Else Find Z in LC Return Z	G4 Unmodified	G5 Unmodified	G6 C(X): If (X,Z) \notin LC $Y2 = H2(X)$ If (X,Y1) \notin L1 $Z \leftarrow \{0,1\}^n$ $Y1 = Z + Y2$ // L1 \leftarrow (X,Y1) : L1 Else Find Y1 $Z = Y1 + Y2$ LC \leftarrow (X,Z) : LC Else Find Z in LC Return Z	G7 C(X): If (X,Z) \notin LC $Y2 = H2(X)$ // If (X,Y1) \notin L1 $Z \leftarrow \{0,1\}^n$ $Y1 = Z + Y2$ // Else Find Y1 // Z = Y1 + Y2 LC \leftarrow (X,Z) : LC Else Find Z in LC Return Z	G8 C(X): If (X,Z) \notin LC // Y2 = H2(X) $Z \leftarrow \{0,1\}^n$ // Y1 = Z + Y2 LC \leftarrow (X,Z) : LC Else Find Z in LC Return Z	G9 Unmodified
H1(X): If (X,Y1) \notin L1 $Y1 \leftarrow \{0,1\}^n$ L1 \leftarrow (X,Y1) : L1 Else Find Y1 on L1 Return Y1	Unmodified	Unmodified	Unmodified	H1(X): If (X,Y1) \notin L1 If (X,Z) \in LC Find Z $Y1 = Z + H2(X)$ Else $Y1 \leftarrow \{0,1\}^n$ $Z = Y1 + H2(X)$ LC \leftarrow (X,Z) : LC L1 \leftarrow (X,Y1) : L1 Else Find Y1 in L1 Return Y1	H1(X): If (X,Y1) \notin L1 If (X,Z) \in LC Find Z // Y1 = Z + H2(X) Else $Z \leftarrow \{0,1\}^n$ LC \leftarrow (X,Z) : LC $Y1 = Z + H2(X)$ L1 \leftarrow (X,Y1) : L1 Else Find Y1 in L1 Return Y1	Unmodified	Unmodified	Unmodified	H1(X): Return C(X) + H2(X)
H2(X): If (X,Y2) \notin L2 $Y2 \leftarrow \{0,1\}^n$ L2 \leftarrow (X,Y2) : L2 Else Find Y2 on L2 Return Y2	Unmodified	Unmodified	Unmodified	Unmodified	Unmodified	Unmodified	Unmodified	Unmodified	Unmodified

G0: The real game.

G1: Keep track of the C queries.

G2: Inline H1.

G3: Change of random variable. Y1 remain iid and random for different X on L1. No change in output distribution.

G4: Two changes: 1) Check if Y1 could have been set in C. This is never reached (since if X in LC it will be automatically on L1).

2) Always add the current entry in L1 to LC (which is necessarily not in LC).

G5: Change of random variable.

G6: Remove handling of L1 from C. Consistency between LC and L1 is now ensured within H1 (due to the **first** change in **G4**).

G7: If X is in L1 it will be also in LC due to the operation of H1. Put differently, if X is not in LC it won't be in L1 either. So the first If check always holds: Remove if-Else.

G8: Y1 and Y2 are not used in C. Defer handling of Y2 to H2 and remove from C.

G9: H1 uses $C(X) + H2(X)$. Irrespectively of (X,Z) being in LC or not this matches the output on **G8** (due to the **second** change in **G4**).

The resulting simulator, which makes at most 1 C-query per H1 call, is:

Sim1^C(X): Return C(X) + Sim2(X)	Sim2(X): If (X,Y2) \notin L2 $Y2 \leftarrow \{0,1\}^n$ L2 \leftarrow (X,Y2) : L2 Else Find Y2 in L2 Return Y2
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